

Chapter 9 DECISIONS, ACTIONS, AND GAMES

While we have now developed separate logics for knowledge update, belief revision, or preference change, concrete instances of rational agency bring in all three entangled. The most concrete setting where this happens is in *games*, and therefore, this chapter is mainly devoted to the question how we can use our earlier logics to describe the latter interactive scenarios. We will be discussing both the logical *statics*, viewing games as unchanging structures representing all possible runs of some process, and the *dynamics* which arises when we make things happen on such a ‘stage’. As we shall see, this involves a mixture of two different faces of logic, or even two different tribal communities, hard-core logics of computation and philosophical logics of belief and preference. But it is precisely this surprising mixture which makes current theories of agency and interaction so lively.

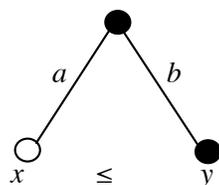
We start with a few examples showing what we are interested in eventually. Then we move to a series of standard logics describing static game structure, from moves to preferences and epistemic uncertainty. More toward the end of the chapter, we introduce the earlier dynamic logics, and see what they have to add in scenarios involving information update and belief revision where given games can change as new information arrives. This chapter is meant to make a connection. It is by no means a full treatment of logical perspectives on games, for which we refer to the companion volume van Benthem, to appear.

9.1 Decisions, practical reasoning, and games

Action and preference Even the simplest scenarios of practical reasoning about agents involve a number of notions at the same time. Recall this example from Chapter 8:

Example One single decision.

An agent has two alternative courses of action, but prefers one outcome to the other:



A proto-typical form of reasoning here would be the ‘Practical Syllogism’:

- (i) the agent can do both a and b ,
 - (ii) the agent prefers the result of a over the result of b ,
- and therefore,
- (iii) the agent will do [or maybe: should do?] b . ■

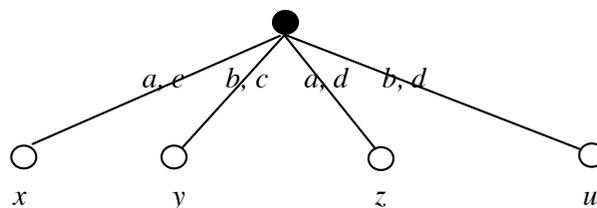
This predictive inference, or maybe requirement, is in fact the basic notion of *rationality* for agents throughout a vast literature in philosophy, economics, and many other fields. It can be used to predict behaviour beforehand, or rationalize observed behaviour afterwards.

Adding beliefs Decision scenarios crucially involve all our notions so far. *Preference* occurs intertwined with *action*, and a reasonable way of taking the conclusion is, not as knowledge ruling out courses of action, but as supporting a *belief* that the agent will take action b : the latter event is now more plausible than the world where she takes action a . Thus, modeling even very simple decision scenarios involves logics of different kinds.

Beliefs come in even more strongly closer to actual decision theory. There, decisions are modeled with uncertainty about possible states of nature, and we are told to choose the action with the highest *expected value*, a probabilistically weighted sum of utility values for the various outcomes. The probability distribution over states of nature represents *beliefs* we have about the world, or the behaviour of an opponent. Here is an yet simpler scenario:

Example Deciding with an external influence.

Nature has two moves c, d , and the agent must now consider combined moves:



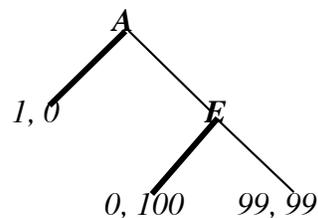
Now, the agent might already have good reasons to think that Nature's move c is more plausible than move d . This will turn the outcomes into an epistemic-doxastic model in the style of Chapter 6: the epistemic range has 4 worlds, but the most plausible ones are just 2: x, y , while the agent's preference might now just refer to the latter area. ■

Thus, we get the entanglement discussed in Chapter 8, while we also see the earlier point in Chapter 6 that it is our beliefs which guide our actions. Here is a typical scenario for this.

Multi-agent decision: ‘solving’ games by Backward Induction In a multi-agent setting, behaviour is locked in place by *mutual expectations*. This requires an interactive decision dynamics, and standard game solution procedures like *Backward Induction* do exactly this sort of thing. Again, recall an earlier example from Chapter 8:

Example Reasoning about interaction.

In the following game tree, players’ preferences are encoded in the utility values, as pairs ‘(value of *A*, value for *E*)’. Backward Induction tells player *E* to turn left when she can, just like in our single decision case, which gives *A* a belief that this would happen, and therefore, based on this belief about his counter-player, *A* should turn left at the start:



Why should players act this way? The reasoning is again a mixture of all notions so far. *A* turns left since she believes that *E* will turn left, and then her preference is for grabbing the value 1. Thus, practical reasoning intertwines action, preference, and belief. ■

Here is the rule which drives all this, at least when preferences are encoded numerically:

Definition Backward Induction algorithm.

Starting from the leaves, one assigns values for each player to each node, using the rule

Suppose *E* is to move at a node, and all values for daughters are known. The *E*-value is the maximum of all the *E*-values on the daughters, the *A*-value is the minimum of the *A*-values at all *E*-best daughters. The dual calculation for *A*'s turns is completely analogous. ■

This rule is so obvious that it never raises objections when taught, and it is easy to apply, telling us what players’ best course of action would be (Osborne & Rubinstein 1994). And yet, it is packed with various assumptions. We will perform a ‘logical deconstruction’ of the underlying reasoning in this chapter, but for now, just note the following features:

- (a) the rule assumes that the situation is viewed in the same way by both players:
since the calculations are totally similar,
- (b) the rule assumes worst-case behaviour on the part of one's opponents,
since we take a minimum of values in case it is not our turn,
- (c) the rule changes its interpretation of the values: at leaves they encode plain utilities, while higher up in the game tree, they represent *expected utilities*.

Thus, in terms of our earlier chapters, we would say that, despite the numerical trappings, Backward Induction is really an inductive mechanism for generating a *plausibility order* among histories, and hence, it relates our models in Chapter 6 to those of Chapter 8: betterness and plausibility order become intertwined in a systematic manner. It is this sort of situation which we should study eventually in the logic of games.

But for now, we step back, and look at what 'logic of games' would involve *ab initio*, starting without any preferences at all. So, let us first consider pure action structure, since even that has a good deal of logic to it, which can be brought out as such.¹⁴⁶ We will add further preferential and epistemic structure toward more realistic games in due course.

9.2 Games and process equivalence

In this chapter, we view extensive games as multi-agent processes which can be studied just like any process in logic and computer science, given the right logical language. Technically, such structures are models for a poly-modal logic in a straightforward sense:

Definition Extensive games.

An *extensive game form* is a tree $M = (NODES, MOVES, \mathit{turn}, \mathit{end}, V)$ which is a modal model with binary transition relations taken from the set *MOVES* pointing from parent to daughter nodes. Also, intermediate nodes have unary proposition letters turn_i indicating the unique player whose turn it is, while *end* marks end nodes without further moves. The valuation *V* for proposition letters may also interpret other relevant predicates at nodes, such as utility values for players or more external properties of game states. ■

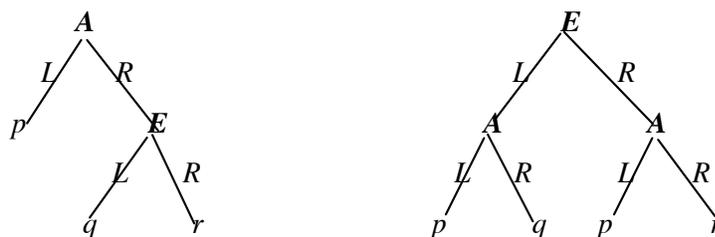
¹⁴⁶ Our treatment will be a sketch: for details see the companion volume van Benthem, to appear.

But do we really just want to jump on board of this analogy, comfortable as it is to a modal logician? Just to emphasize how close, indeed, this setting is to our earlier style of analysis, here is a fundamental issue of invariance we raised before in Chapter 2. At which level do we want to operate in the logical study of games, or in more Clintonesque terms:

When are two games are the same?

Example The same game, or not?

As a simple example that is easy to remember, consider the following two games:



Are these the same? As with general processes in computer science (cf. Chapter 2), the answer crucially depends on our level of interest in the details of what is going on:

(a) *If we focus on turns and moves, then the two games are not equivalent.*

For they differ in ‘protocol’ (who gets to play first) and in choice structure. For instance, the first game, but not the second has a stage where it is up to *E* to determine whether the outcome is *q* or *r*. This is indeed a natural level for looking at game, involving local *actions* and choices, as encoded in modal *bisimulations* – and the appropriate language will be a standard modal one. But one might also want to call these games equivalent in another sense: looking at achievable outcomes only, and players powers for controlling these:

(b) *If we focus on outcome powers only, then the two games are equivalent.*

The reason is that, regardless of protocol and local choices, players can force the same sets of eventual outcomes across these games, using strategies that are available to them:

A can force the outcome to fall in the sets $\{p\}, \{q, r\}$,

E can force the outcome to fall in the sets $\{p, q\}, \{p, r\}$.

In the left-hand tree, *A* has 2 strategies, and so does *E*, yielding the listed sets. In the right-hand tree, *E* has 2 strategies, while *A* has 4: *LL*, *LR*, *RL* and *RR*. Of these, *LL* yields the

outcome set $\{p\}$, and RR yields $\{q, r\}$. But LR, RL guarantee only supersets $\{p, r\}, \{q, p\}$ of $\{p\}$: i.e., weaker powers. Thus the same 'control' results in both games. ■

In this book, we will concentrate on the extensive modal level, but we do note that the coarser power level is natural, too. It is closer to the 'strategic forms' for games employed in game theory, and also, it fits very well with 'logic games' (cf. our companion volume van Benthem, to appear). We take our leave of the latter connection with one illustration:

Remark: game equivalence as logical equivalence In an obvious sense, the two games in the preceding example represent the two sides of the following valid logical law

$$p \wedge (q \vee r) \leftrightarrow (p \wedge q) \vee (p \wedge r) \quad \text{Distribution}$$

Just read conjunction and disjunction as choices for different players. In a global input–output view, Distribution switches scheduling order without affecting players' powers. ■

9.3 Basic modal action logic of extensive games

Henceforth, we will restrict attention to extensive games in the modal style.

Basic modal logic On extensive game trees, a standard modal language works as follows:

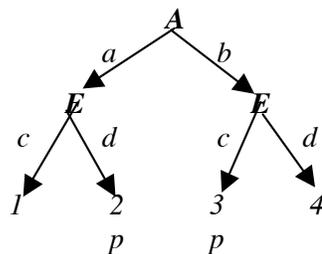
Definition Modal game language and semantics.

Modal formulas are interpreted at nodes s in game trees \mathcal{M} . Labeled modalities $\langle a \rangle \phi$ express that some move a is available leading to a next node in the game tree satisfying ϕ . Proposition letters true at nodes may be special-purpose constants for game structure, such as indications for turns and end-points, but also arbitrary local properties. ■

In particular, modal operator combinations now describe potential interaction:

Example Modal operators and strategic powers.

Consider a simple 2-step game like the following, between two players A, E :



Player E clearly has a strategy making sure that a state is reached where p holds. And this feature of the game is directly expressed by the modal formula $[a]\langle d\rangle p \wedge [b]\langle c\rangle p$. ■

More generally, letting *move* stand for the union of all relations available to players, in the preceding game, the modal operator combination

$$[\text{move-A}]\langle \text{move-E} \rangle \phi$$

says that, at the current node of evaluation, player E has a strategy for responding to A 's initial move which ensures that the property expressed by ϕ results after two steps of play.

¹⁴⁷ This fundamental relation between strategic structure and logical operator combinations is crucial to so-called ‘*logic games*’, a topic for which we refer to van Benthem 2007.

Excluded middle and determinacy Extending this observation to extensive games up to some finite depth k , and using alternations $\square\Diamond\square\Diamond\dots$ of modal operators up to length k , we can express the existence of winning strategies in fixed finite games. Indeed, given this connection, with finite depth, standard logical laws have immediate game-theoretic import. In particular, consider the valid *law of excluded middle* in the following modal form

$$\square\Diamond\square\Diamond\dots\phi \vee \neg\square\Diamond\square\Diamond\dots\phi$$

or after some logical equivalences, pushing the negation inside:

$$\square\Diamond\square\Diamond\dots\phi \vee \Diamond\square\Diamond\square\dots\neg\phi,$$

where the dots indicate the depth of the tree. Here is its game-theoretic import:

Fact Modal excluded middle expresses the determinacy of finite games.

Here, *determinacy* is the fundamental property of many games that *one of the two players has a winning strategy*. This need not be true in infinite games (players cannot both have one, but maybe neither has), and descriptive set theory has deep theory around this notion.

Zermelo's theorem This brings us to perhaps the oldest game-theoretic result proved in mathematics, even predating Backward Induction, proved by Ernst Zermelo in 1913:

¹⁴⁷ One can also express the existence of ‘winning strategies’, ‘losing strategies’, and so forth.

Theorem Every finite zero-sum 2-player game is determined.

Proof Here is a simple algorithm determining the player having the winning strategy at any given node of a game tree of this finite sort. It works bottom-up through the game tree. First, colour those end nodes *black* that are wins for player *A*, and colour the other end nodes *white*, being the wins for *E*. Then extend this colouring stepwise as follows:

If all children of node *s* have been coloured already, do one of the following:

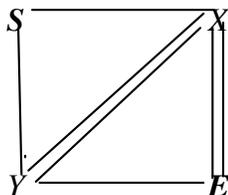
- (a) if player *A* is to move, and at least one child is black:
colour *s* *black*; if all children are white, colour *s* *white*
- (b) if player *E* is to move, and at least one child is white:
colour *s* *white*; if all children are black, colour *s* *black*

This procedure eventually colours all nodes black where player *A* has a winning strategy, while making those where *E* has a winning strategy white. And the reason for its correctness is that a player has a winning strategy at one of her turns iff she can make a move to at least one daughter node where she has a winning strategy. ■

Zermelo's Theorem is widely applicable. Recall the Teaching Game from our Introduction:

Example Teaching, the grim realities.

A Student located at position *S* in the next diagram wants to reach the *escape* *E* below, while the Teacher wants to prevent him from getting there. Each line segment is a path that can be traveled. In each round of the game, the Teacher cuts one connection, anywhere, while the Student can, and must travel one link still open to him at his current position:



General education games like this arise on any graph with single or multiple lines. ■

We now have a principled explanation why either Student or Teacher has a winning strategy, since this game is obviously two-player zero sum and of finite depth – though there need be an effective computational method for solving the game. Zermelo's Theorem

implies that in Chess, one player has a winning strategy, or the other a non-losing one – but almost a century later, we do not know which is which: the game tree is just too large.¹⁴⁸

9.4 Fixed-point languages for equilibrium notions

A good test for logical languages is checking their expressive power for rendering proofs of significant results. Here our basic modal language seems a bit poor, since it cannot express the *generic character* of the Zermelo solution. The following is what the colouring algorithm really says. Starting from atomic predicates win_i at end nodes indicating which player has won, we inductively defined predicates WIN_i ('player i has a winning strategy at the current node') through the following recursion:

$$\mathit{WIN}_i \leftrightarrow (\mathit{end} \ \& \ \mathit{win}_i) \vee (\mathit{turn}_i \ \& \ \langle E \rangle \mathit{WIN}_i) \vee (\mathit{turn}_j \ \& \ [A] \mathit{WIN}_i)$$

Here E is the union of all available moves for player i , and A that of all moves for the counter-player j . This schema is an inductive definition for the predicate WIN_i , which we can also write as a *smallest fixed-point* expression in an extended modal language:

Fact The Zermelo solution is definable as follows in the modal μ -calculus:

$$\mathit{WIN}_i = \mu p \bullet (\mathit{end} \ \& \ \mathit{win}_i) \vee (\mathit{turn}_i \ \& \ \langle E \rangle p) \vee (\mathit{turn}_j \ \& \ [A] p)^{149}$$

Here the formula on the right-hand side belongs to the *modal μ -calculus*, an extension of the basic modal language with operators for *smallest* (and *greatest*) *fixed-points* defining inductive notions. This system was originally invented to increase the power of modal logic as a process theory. We refer to the literature for details, cf. Bradfield & Stirling 2007, van Benthem 2008. For here, we just observe that fixed-points fit well with strategic equilibria, and the μ -calculus has many further uses in games (cf. Chapter 14 for another example).

As we said already in the global perspective of Section 9.2, games are all about powers of control which players have over 'reachable propositions' via their strategies:

¹⁴⁸ But the clock is ticking for Chess. Recently, for the game of *Checkers*, 15 years of real-time computer verification yielded the Zermelo answer: the starting player has a non-losing strategy.

¹⁴⁹ Note that the defining schema only has *syntactically positive occurrences* of the predicate p .

Definition Forcing modalities.

Forcing modalities are interpreted as follows in extensive game models as defined earlier: $\mathbf{M}, s \models \{i\}\phi$ iff player i has a strategy for the sub-game starting at s which guarantees that only nodes will be visited where ϕ holds, whatever the other player does. ■

Forcing talk is widespread in games, and it is an obvious target for logical formalization:¹⁵⁰

Fact The modal μ -calculus can define forcing modalities.

Proof The formula $\{i\}\phi = \mu p \bullet (\mathbf{end} \ \& \ \phi) \vee (\mathbf{turn}_i \ \& \ \langle E \rangle p) \vee (\mathbf{turn}_j \ \& \ [A]p)$ defines the existence of a strategy for i ensuring that proposition ϕ holds, whatever the other plays. ■

But many other notions are definable. For instance, the recursion

$$\mathbf{COOP} \ \phi \Leftrightarrow \mu p \bullet (\mathbf{end} \ \& \ \phi) \vee (\mathbf{turn}_i \ \& \ \langle E \rangle p) \vee (\mathbf{turn}_j \ \& \ \langle A \rangle p)$$

defines the existence of a *cooperative outcome* ϕ , just by shifting some modalities.¹⁵¹

Digression: from smallest to greatest fixed-points The above modal fixed-point definitions reflect the equilibrium character of basic game-theoretic notions (Osborne & Rubinstein 1994), reached through some process of iteration. In this general setting, which includes infinite games, we would switch from smallest to *greatest fixed-points*, as in the formula

$$\{i\}\phi = \nu q \bullet (\phi \ \& \ (\mathbf{turn}_i \ \& \ \langle \mathbf{move}\text{-}i \rangle q) \vee (\mathbf{turn}_j \ \& \ [\mathbf{move}\text{-}j]q)).$$

This is also more in line with our intuitive view of strategies. The point is not that they are built up from below, but that they can be used as needed, and then remain at our service as pristine as ever the next time – the way we think of doctors. This is the modern perspective of *co-algebra* (Venema 2007). More generally, greatest fixed-points seem the best logical analogue to the standard equilibrium theorems from analysis that are used in game theory.

¹⁵⁰ Note that $\{i\}\phi$ talks about intermediate nodes, not just end nodes of a game. The existence of a winning strategy for player i can then be formulated by restricting to endpoints: $\{i\}(\mathbf{end} \rightarrow \mathbf{win}_i)$.

¹⁵¹ This fixed point can still be defined in propositional dynamic logic, using the formula $\langle ((\mathbf{turn}_i)? ; E) \cup (\mathbf{turn}_j)? ; A \rangle^* (\mathbf{end} \ \& \ \phi)$, – but we will only use the latter system later in the game setting.

But why logic? This may be a good place to ask the question *what is the point* of giving logical definitions of game-theoretic notions? The way I see this chapter, logic has the same virtues for games as it has elsewhere. Through formalization of a given practice, we see more clearly what makes its key notions tick, and we also get a feel for new alternative notions, as the logical language has myriads of possible definitions. Moreover, the theory of expressive power, completeness, and complexity of our logics can be used for model checking, proof search, and other activities not normally associated with game theory.

But there is also another viewpoint: not central to this book on information dynamics, but equally important in general. As we saw in Chapter 2, basic notions of logic *themselves* have a game character, such as argumentation, model checking, or model comparison. Thus, logic does not just *describe games*, it also *embodies games*. When we pursue the interface in this dual manner, the true grip of the logic & games connection becomes clear, and this will be the main thrust of the companion volume van Benthem, to appear.

9.5 Dynamic logics of strategies

Given that strategies, rather than single moves, are obvious protagonists in our story, it makes sense to move them into our logics. This requires another extension of modal logic, viz. *propositional dynamic logic* (*PDL*, Bradfield & Stirling 2007) used for analyzing conversation in Chapter 3, and for common knowledge in Chapter 5. *PDL* was created to describe the structure and effects of imperative computer programs constructed using the sequential operations of (a) sequential composition $;$, (b) guarded choice *IF ... THEN... ELSE...*, and (c) guarded iterations *WHILE... DO...* We recall some earlier definitions:

Definition Propositional dynamic logic.

The *language of PDL* defines formulas and programs in a mutual recursion, with formulas denoting sets of worlds (‘local conditions’ on ‘states’ of the process), while programs denote binary transition relations between worlds, recording pairs of input and output states for their successful terminating computations. Programs are created from

atomic actions (‘moves’) a, b, \dots and tests $?\varphi$ for arbitrary formulas φ ,
 using the three operations of $;$ (interpreted as sequential composition),
 \cup (non-deterministic choice) and $*$ (non-deterministic finite iteration).

Formulas are as in our basic modal language, but with modalities $[\pi]\varphi$ saying that φ is true after every successful execution of the program π starting at the current world. ■

The logic *PDL* is decidable, and it has an transparent complete set of axioms for describing validity which we will give later. Right here, our point is that this formalism can say a lot more about our preceding examples. For instance, the *move* relation in our discussion of our first extensive game in this chapter was really a union of atomic transition relations, and the pattern that we discussed for the winning strategy was as follows:

$$[a \cup b] \langle c \cup d \rangle p.$$

Strategies as transition relations But we can deal with strategies much more explicitly. Game-theoretic strategies may be viewed as *partial transition functions* defined on players' turns, given via a bunch of conditional instructions of the form "if she plays this, then I play that". More generally, strategies may be viewed as binary transition relations, allowing for non-determinism, i.e., more than one 'optimal move'. While this is not the standard approach in game theory, such a relational view comes much closer to *plans* that agents have in interactive settings. A plan can be extremely useful, even when it merely constrains my future moves, without necessarily fixing my course of action. Thus, on top of the 'hard-wired' move relations in a game, we now get defined further relations, corresponding to players' strategies, and these definitions can often be given explicitly in a *PDL*-format.

In particular, in finite games, we can define an explicit version of the earlier forcing modality, indicating the strategy involved – without recourse to the modal μ -calculus:

Fact For any game program expression σ , *PDL* can define an explicit forcing modality $\{\sigma, i\}\phi$ stating that σ is a strategy for player i forcing the game, against any play of the others, to pass only through states satisfying ϕ .

The precise definition is an easy exercise (cf. van Benthem 2002). Also, given strategies for both players, we should get to a unique history of a game, and here is how:

Fact Outcomes of running joint strategies σ, τ can be defined in *PDL*.

Proof The formula $[((?turn_E; \sigma) \cup (?turn_A; \tau))^*] (end \rightarrow p)$ does the job. ¹⁵² ■

But also, on a model-by-model basis, the expressive power of *PDL* for defining specific strategies is high. Consider any finite game \mathbf{M} with a strategy σ for player i . As a relation, σ is a finite set of ordered pairs (s, t) . Thus, it can be defined by enumeration as a program union, if we define these ordered pairs. To do so, assume we have an ‘expressive’ model \mathbf{M} , where states s are definable in our modal language by formulas def_s . ¹⁵³ Then we define transitions (s, t) by formulas $def_s; a; def_t$, with a being the relevant move:

Fact In expressive finite extensive games, all strategies are *PDL*-definable.

Of course, this is a trivial result, but it does suggest that *PDL* is on the right track. Dynamic logic can also define strategies running over only part of a game, and their *combination*. The following modal operator describes the effect of such a partial strategy σ for player E running until the first game states where it is no longer defined:

$$\{\sigma, E\}\phi \quad [(?turn_E; \sigma) \cup (?turn_A; move-A)^*] \phi$$

A basic operation on partial strategies would be *intersection*, leaving only moves allowed by both. Van Benthem 2002 discusses Relational Algebra as a form of strategy calculus. Van Benthem 2007 is a full-fledged survey and defense of modal logics with explicit strategies. Indeed, propositional dynamic logic does a reasonable job in defining explicit strategies in simple extensive games. ¹⁵⁴ In the next sections, we will extend it to deal with more realistic game structures, such as preferences and imperfect information.

¹⁵² Dropping the antecedent ‘*end* \rightarrow ’ here will describe effects of strategies at intermediate nodes.

¹⁵³ This expressive power can be achieved in several ways: e.g., using temporal *past modalities* over converse moves which can describe the total history leading up to s .

¹⁵⁴ *Stronger modal logics of strategies?* The modal μ -calculus is a natural extension of *PDL*, but it lacks explicit programs or strategies, as its formulas merely define properties of states. Is there a version of the μ -calculus which extends *PDL* in defining more transition relations? Say, a simple strategy ‘keep playing a ’ guarantees infinite a -branches for greatest fixed-point formulas like $\nu p \bullet \langle a \rangle p$. Van Benthem & Ikegami 2008 look at richer fragments than *PDL* with explicit programs as solutions to fixed-point equations of special forms, guaranteeing uniform convergence by stage ω .

9.6 Preference logic and defining backward induction

Real games, as we have seen several times already, go beyond mere game forms by adding preferences for players over outcome states, or numerical utilities beyond ‘win’ and ‘lose’. In this area, defining the earlier-discussed Backward Induction procedure for solving extensive games, rather than computing binary Zermelo winning positions, has become a benchmark for game logics – and many solutions exist:

Fact The Backward Induction path is definable in modal preference logic.

Solutions have been published by many logicians and game-theorists in recent years, cf. de Bruin 2004, van der Hoek & Pauly 2007. We do not state an explicit *PDL*-style solution here, but we give one version involving the modal preference language of Chapter 8:

$\langle \text{pref}_i \rangle \phi$: player i prefers some node where ϕ holds to the current one.

The following result from van Benthem, van Otterloo & Roy 2006 defines the backward induction path as a unique relation σ : not by means of any specific modal formula in game models \mathbf{M} , but rather via the following *frame correspondence* on finite structures:

Fact The *BI* strategy is definable as the unique relation σ satisfying the following axiom for all propositions P – viewed as sets of nodes –, for all players i :

$(\text{turn}_i \ \& \ \langle \sigma^* \rangle (\text{end} \ \& \ P)) \rightarrow [\text{move-}i] \langle \sigma^* \rangle (\text{end} \ \& \ \langle \text{pref}_i \rangle P)$.

Proof The axiom expresses a form of rationality: at the current node, no alternative move for a player guarantees outcomes that are all strictly better than the outcomes ensuing from playing the current backward induction move. The proof of this frame correspondence is a straightforward induction on levels of the finite game tree. ■

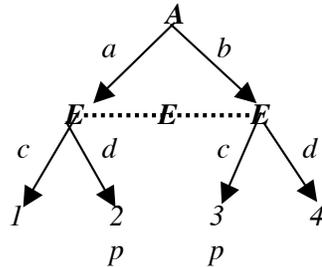
9.7 Epistemic logic of games with imperfect information

The next level of static structure gives up another presupposition made in this chapter so far, that of perfect information for players concerning their situation. Consider extensive games with *imperfect information*, whose players need not know where they are in the game tree. This happens in card games, electronic communication, and real life game scenarios with bounds on memory or observation. Such games have ‘information sets’: equivalence classes of epistemic relations \sim_i between nodes which players i cannot distinguish, as in the

S5-type models of Chapters 2–5. Van Benthem 2001 points out how these games model a combined epistemic modal language including knowledge operators $K_i\phi$ interpreted in the usual manner as " ϕ is true at all nodes \sim_i -related to the current one".

Example Partial observation in games.

In this imperfect information game, a dotted line indicates player E 's uncertainty about her position when her turn comes. Thus, she does not know the move played by player A :¹⁵⁵



Structures like this are game models of the earlier kind with added epistemic *uncertainty relations* \sim_i for each player. Thus, they interpret a combined *dynamic-epistemic* language. For instance, after A plays move c in the root, in both middle states, E knows that playing a or b will give her p – as the disjunction $\langle a \rangle p \vee \langle b \rangle p$ is true at both middle states:

$$K_E(\langle a \rangle p \vee \langle b \rangle p)$$

On the other hand, there is no *specific* move of which E knows at this stage that it will guarantee a p -outcome – and this shows in the truth of the formula

$$\neg K_E \langle a \rangle p \ \& \ \neg K_E \langle b \rangle p$$

Thus, E knows *de dicto* that she has a strategy which guarantees p , but she does not know, *de re*, of any specific strategy that it guarantees p . Such finer distinctions are typical for a language with both actions and knowledge for agents.¹⁵⁶

We can analyze imperfect information games studying properties of players by modal frame correspondences. An example is an earlier analysis of Perfect Recall for a player i :

¹⁵⁵ Maybe A put his move in an envelope, or E was otherwise prevented from observing.

¹⁵⁶ You may know that the ideal partner for you is around on the streets, but tragically, you might never convert this $K\exists$ combination into $\exists K$ knowledge that some particular person is right for you.

Fact The axiom $K_i[a]\phi \rightarrow [a]K_i\phi$ holds for player i w.r.t. any proposition ϕ iff \mathbf{M} satisfies *Confluence*: $\forall xyz: ((xR_a y \ \& \ y \sim_i z) \rightarrow \exists u: ((x \sim_i u \ \& \ u R_a z))$.

Similar frame analyses work for memory bounds, and observational powers for agents (van Benthem 2001). For instance, in the opposite direction of Perfect Recall, agents satisfy the principle of ‘No Miracles’ when their epistemic uncertainty relations can only disappear when they observe two subsequent events which they can distinguish.¹⁵⁷ Incidentally, the game which we just gave does satisfy Perfect Recall, but it violates No Miracles, since \mathbf{E} suddenly knows exactly where she is after she has played her move. We will discuss these properties in much greater generality in the epistemic-temporal logics of Chapter 11.

Uniform strategies Another striking aspect of the above game with imperfect information is its *non-determinacy*. \mathbf{E} ’s playing ‘the opposite direction from that chosen by player \mathbf{A} ’ was a strategy guaranteeing outcome p in the original game – but it is unusable now. For, \mathbf{E} cannot tell if the condition holds! Game theorists only accept *uniform strategies* here, prescribing the same move at indistinguishable nodes. But then no player has a winning strategy, when we interpret p as ‘ \mathbf{E} wins’ (and hence $\neg p$ as a win for player \mathbf{A}). \mathbf{A} did not have one to begin with, \mathbf{E} loses hers. The game does have probabilistic solutions in random strategies:¹⁵⁸ essentially, it is like Matching Pennies, but we do not pursue these here.

Now once again for explicit strategies! As before, we can add *PDL*-style programs here to define players’ strategies under the new circumstances. But there is a twist. Especially relevant then are the ‘*knowledge programs*’ of Fagin et al. 1995, whose only test conditions for actions are knowledge statements for agents. In such programs, the actions prescribed for an agent can only be guarded by conditions which the agent knows to be true or false. It is easy to see that knowledge programs can only define *uniform strategies*, i.e., transition relations where a player always chooses the same move at any two game nodes which she cannot distinguish epistemically. A converse also holds, modulo some assumptions on expressiveness of the game language defining nodes in the game tree (van Benthem 2001):

¹⁵⁷ Our analysis restates that of Halpern & Vardi on epistemic-temporal logic (Fagin et al. 1995).

¹⁵⁸ Both players should play both moves with probability $\frac{1}{2}$, for an optimal outcome 0 .

Fact On expressive finite games of imperfect information, the uniform strategies are precisely those definable by knowledge programs in epistemic *PDL*.

9.8 From statics to dynamics: *DEL*-representable games

Now we make a switch. We have linked logic to games in various ways, but our approach was *static* in the sense of Chapter 1, using modal-preferential-epistemic logics to describe properties of given games, which themselves do not change. While this is in line with the literature in the area, it also makes sense to look at *dynamic* scenarios in the sense of this book, where games themselves can change because of certain triggering events. But as an intermediate step toward that goal, we now first analyze how a static game model might have come about as the result of some dynamic process – the way we see a dormant volcano but can also imagine the tectonic forces that shaped it originally. The present and the next section provide two illustrations of this perspective, linking games of imperfect information first to *DEL* (Chapter 4), and then to epistemic-temporal logics (Chapter 11).

Imperfect information games and dynamic-epistemic logic The reader should long have seen the connection between imperfect information games and the product update system *DEL* of our Chapter 4, as it creates successive models starting from some initial one. For a start, as with adding preferences, there are *two levels* for making the earlier action logic *PDL* epistemic. One can connect worlds, as with the above language with standard epistemic modalities K_i . But one can also place epistemic structure *on the moves themselves*, as in dynamic epistemic logic. This at once raises the issue which imperfect information games ‘make sense’ in terms of real scenarios from dynamic-epistemic logic, as opposed to being just arbitrary placements of uncertainty links over game forms:

What sort of imperfect information games correspond to the update evolution of some initial model using product update as in dynamic-epistemic logic?

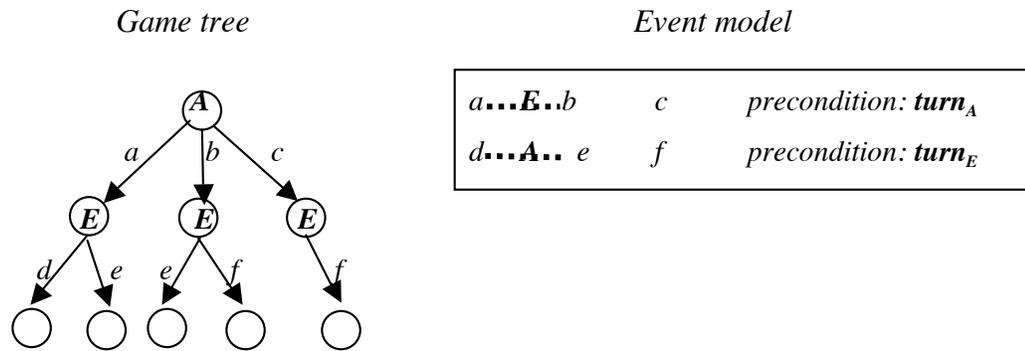
Theorem An extensive game tree is isomorphic to a repeated product update model $Tree(M, E)$ over some epistemic event model E iff it satisfies, for all players:

- (a) Perfect Recall,
- (b) (uniform) No Miracles,

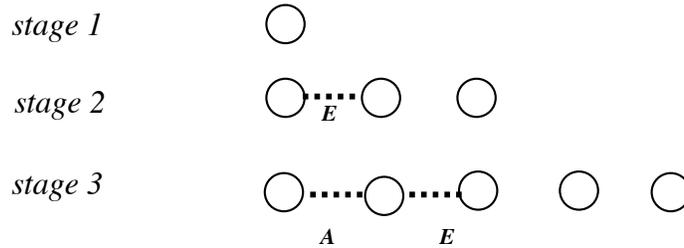
(c) Bisimulation Invariance for the domains of the move relations.¹⁵⁹

Here Perfect Recall is essentially the earlier commutation property between moves and uncertainty. We do not prove the Theorem here, as an extensive version will be presented in Chapter 11 below (van Benthem & Liu 2004, van Benthem, Gerbrandy & Pacuit 2007 have definitions and proofs). Instead, here is an illustration. We can create an imperfect information game for *DEL*-style players by first specifying preconditions and observation structure on the moves, and then producing the game tree through iterated product update:

Example Updates during play: propagating ignorance along a game tree.



Here are the successive updates that create the right uncertainty links:



9.9 Future uncertainty, procedural information, and branching temporal logic

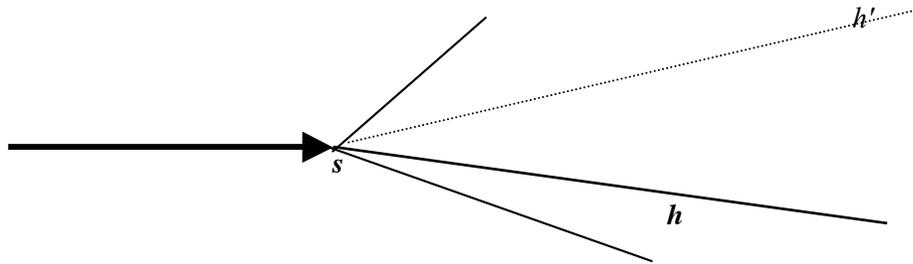
A second logical perspective on games with imperfect information starts from the following observation. The phrase ‘imperfect information’ covers two intuitively different senses. One is *observation uncertainty*: players may not have been able to observe all events so far, and so they do not know where they are in the game tree. This is the ‘past-oriented’ view of *DEL*. But there is also ‘future-oriented’ *expectation uncertainty*, which

¹⁵⁹ I.e., two epistemically bisimilar nodes in the game tree make the same moves executable.

occurs even in games of ‘perfect’ information players may know precisely where they are, but they do not know where the game is heading as they do not know what others, or they themselves, are going to do. The positive side of the situation is this. In general, players will have a certain amount of *procedural information* about what is going to happen, and van Benthem 2008 has argued that this is a fundamental and new notion of information concerning the current process agents are in, which may be considered *sui generis*.

But whether viewed negatively or positively, the latter future-oriented kind of knowledge and ignorance need not be reducible to the earlier uncertainty between local nodes. Instead, it naturally suggests current uncertainty between whole future histories, or between players’ strategies (i.e., whole ways in which the game might evolve). The next points are from van Benthem 2004 on information update and belief revision in game trees.

Branching epistemic temporal models The following structure is common to many fields in computer science and philosophy (cf. Chapter 11 below for more details). In tree-like models for branching time, ‘legal histories’ h represent possible evolutions of a given game. At each stage of the game, players are in a node s on some actual history whose past they know, either completely or partially, but whose future is yet to be fully revealed:



This can be described in an action language with knowledge, belief, and added temporal operators. We first describe games of perfect information (about the past, that is):

- (a) $\mathbf{M}, h, s \models F_a \phi$ iff $s \frown \langle a \rangle$ lies on h and $\mathbf{M}, h, s \frown \langle a \rangle \models \phi$
- (b) $\mathbf{M}, h, s \models P_a \phi$ iff $s = s' \frown \langle a \rangle$, and $\mathbf{M}, h, s' \models \phi$
- (c) $\mathbf{M}, h, s \models \diamond_i \phi$ iff $\mathbf{M}, h', s \models \phi$ for some h' equal for i to h up to stage s .

Now, as moves are played publicly, players make public observations of them:

Fact The following valid principle is the temporal equivalent of the key *DEL* reduction axiom for public announcement: $F_a \diamond \phi \leftrightarrow (F_a T \ \& \ \diamond F_a \phi)$.¹⁶⁰

We will see this principle return as a recursion law for *PAL* with *protocols* in Chapter 11.

Trading future for current uncertainty in games Again, there is a possible ‘dynamic reconstruction’ for these models, this time, providing an alternative description closer to the local dynamics of *PAL* and *DEL* after all. Intuitively, each move by a player is a public announcement that changes the current game. Here is a folklore observation converting ‘global’ uncertainty about the future into ‘local’ uncertainty about the present:

Fact Trees with future uncertainty are isomorphic to trees with current uncertainties.

Proof The construction is this (cf. van Benthem 2004 for a slightly more formal version): Given any game tree G , assign epistemic models M_s to each node s whose domain is the set of complete histories passing through s (which all share the same past up to s), letting the agent be uncertain about all of them. ‘Worlds’ in these models may be seen as pairs (h, s) with h any history passing through s . This will cut down the current set of histories in just the right manner. The above epistemic-temporal language matches this construction. ■

9.10 Intermezzo: three levels of logical game analysis

At this point, it may be useful to distinguish three natural levels at which games have given rise to models for logics. All three come with their own intuitions, both static and dynamic.

Level One is the most straightforward perspective: *extensive game trees* themselves are models for modal languages, with nodes as worlds, and accessibility relations over these for actions, preferences, and uncertainty. *Level Two* rather looks at extensive games as *branching tree models*, involving both nodes and complete histories, supporting richer epistemic-temporal (-preferential) languages. The difference with Level One may seem slight in finite extensive games, where histories may be marked by their end-points, so that a point on a given history may also be viewed as just a pair of nodes. But the intuitive step

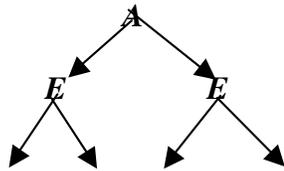
¹⁶⁰ As in our earlier modal-epistemic analysis, commutation of a temporal and an epistemic operator implies a form of *Perfect Recall*: agents’ present uncertainties are always inherited from past ones.

from the first to the second perspective is clear, and also, Level Two cannot be reduced in this manner when game trees are infinite. But even this is not enough for some purposes!

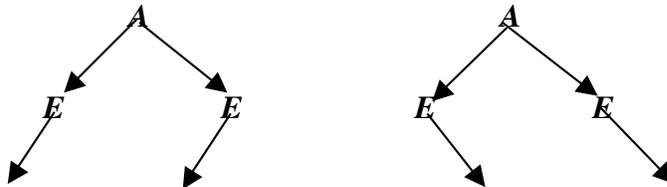
Consider ‘higher’ hypotheses about the future, involving procedural information about other players’ strategies. I may know that I am playing against either a ‘simple automaton’, or a ‘sophisticated learner’. Modeling this may go beyond epistemic-temporal models:

Example Strategic uncertainty.

In the following simple game, let A know that E will play the same move throughout:



Then all four histories are still possible. But A only considers two *future trees* possible, viz.



In longer games, this difference in modeling can be highly important, because observing only one move by E will tell A exactly what E 's strategy will be in the whole game. ■

To model these richer settings, one needs full-fledged *Level Three* epistemic game models.

Definition Epistemic game models.

Epistemic game models for an extensive game G are epistemic models $M = (W, \sim_i, V)$ whose worlds are abstract indices including local (factual) information about all nodes in G , plus whole strategy profiles for players, i.e., total specifications of everyone's behaviour throughout the game. Players' global information about game structure and procedure is encoded by uncertainty relations \sim_i between the worlds of the game model. ■

The above uncertainty between two strategies of my opponent would be naturally encoded in constraints on the set of strategy profiles represented in such a model. And observing

some moves of yours in the game telling me which strategy you are actually following then corresponds to dynamic update of the initial model, in the sense of our earlier chapters.

While we will not use Level-Three models much in this chapter, they are the natural limit for games and other scenarios of interactive agency. Throughout this chapter, we follow the policy of discussing issues at the simplest modeling level where they make sense.

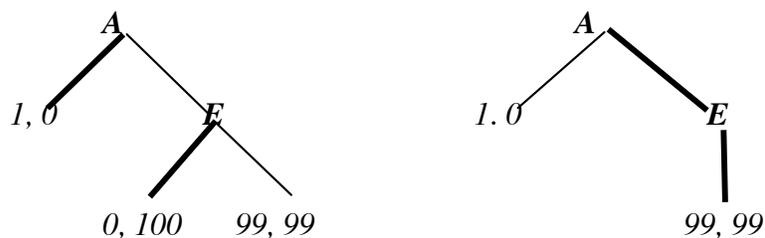
9.11 Game change: public announcements, promises and solving games

Now let us look at actual transformations of given games, and explicit triggers for them.

Promises and intentions Following van Benthem 2007, one can break the impasse of a bad Backward Induction solution by changing the game through making suitable *promises*.

Example Promises and game change.

In the following game, discussed several times before, the ‘bad Nash equilibrium’ $(1, 0)$ can be avoided by *E*’s *promise* that she will not go left, which may be modeled as a public announcement that some histories will not occur (we may make this binding in a number of ways, e.g., by attaching a huge fine to infractions) – and the new equilibrium $(99, 99)$ results, making both players better off by restricting the freedom of one of them!



But one can also add new moves to a game. ¹⁶¹ Van Otterloo 2005 has a dynamic logic of strategic ‘enforceability’ plus preference, where game models can change by announcing players’ intentions. Related scenarios let public announcements provide information about players’ preferences. Our methods from Chapters 3 and following apply here at once:

¹⁶¹ Yes, in this way, one could code up all such game changes beforehand in one grand initial ‘Super Game’ – but that would lose all the flavour of understanding what happens in a stepwise manner.

Theorem Modal logic of games plus public announcement is completely axiomatized by the modal game logic chosen, the recursion axioms of *PAL* for atoms and Booleans, plus the following law for the move modality:

$$\langle !P \rangle \langle a \rangle \varphi \Leftrightarrow (P \wedge \langle a \rangle (P \wedge \langle !P \rangle \varphi)).$$

Strategies once more We can also talk about explicit strategies in all of this. Using *PDL* for strategies in games as before, this leads to the obvious logic *PDL+PAL* adding public announcements $[!P]$. It is easy to show that *PDL* is closed under relativization to definable sub-models, both in its propositional and its program parts, using a recursive operation π/P for programs π which surrounds every atomic move with tests $?P$.

Theorem *PDL+PAL* is axiomatized by merging their separate laws while adding the following reduction axiom: $[!P]\{\sigma\}\phi \Leftrightarrow (P \rightarrow \{\sigma/P\}[!P]\phi)$.¹⁶²

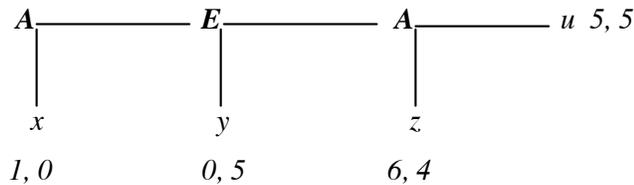
But of course, we also want to know about versions with epistemic preference languages – and hence there are many further questions following up on these initial observations.

Solving games by announcements of rationality Another type of application of public announcement in games iterates assertions expressing that players are rational, as a sort of ‘public reminders’. Van Benthem 2003 has the following result for extensive games:

Theorem The Backward Induction solution for extensive games is obtained through repeated announcement of the assertion “no player chooses a move all of whose further histories end worse than all histories after some other available move”.

Proof This can be proved by a simple induction on finite game trees. The principle behind the argument will be clear by seeing how the announcement procedure works for the following simple ‘Centipede game’, with three turns as indicated, four branches indicated by name, and the pay-offs given for *A*, *E* in that order:

¹⁶² This axiom says what old plan I should have had to make some plan work in the new game. But usually, I already have a plan σ for playing G to get effect ϕ . Now G changes to G' : my game got more, or fewer, moves. How should I revise that current plan σ to get some intended effect ψ in the new game G' ? This may be related to open model-theoretic problems like finding a Los-Tarski theorem syntactically characterizing those *PDL* formulas which are preserved under sub-models.



Stage 0 of the announcement procedure rules out branch u ,

Stage 1 then rules out z , while Stage 2 finally rules out y . ■

As in the analysis of iterated assertions in Chapter 3, this iterated announcement procedure for extensive games (or for the temporal models of the preceding section) ends in largest sub-models in which players have common belief of rationality, or other assertions.¹⁶³

Of course, our logical language provides many further assertions that can be announced, far beyond the ruthless egotism of Backward Induction. Van Benthem 2003, 2004 discuss history-oriented alternatives, where players steer future actions by reminding themselves of *legitimate rights of other players*, because of ‘past favours received’.^{164 165}

The same ideas work in strategic games, with epistemic models M of strategy profiles with preferences and uncertainties for players who know their own strategy, but not that of the others. In Chapter 14, we formulate assertions of *Weak Rationality* (“no player chooses a move which she knows to be worse than some other available one”) and *Strong Rationality* (“each player chooses a move she thinks may be the best possible one”). When announced, these assertions eliminate worlds, and iterating these announcements, perhaps infinitely, one eventually gets to a smallest sub-model where WR or SR are common knowledge:

¹⁶³ What about an epistemic-doxastic temporal preference language where the final sub-models can be *explicitly defined*? Does it need fixed-point operators, like for strategic games in Chapter 14?

¹⁶⁴ Type spaces for games, too, allow for much greater freedom in assumptions about players.

¹⁶⁵ The different options are somewhat reminiscent of a distinction in social philosophy, between Rawls’ forward-oriented *adjustment view* of fairness which recommends minimizing differences in the current income distribution, versus Nozicks’ historical *entitlement view*, which says that a current income distribution is fair if there was a history of fair transactions leading to it.

Theorem The result of iterated announcement of *WR* is the usual solution concept of Iterated Removal of Strictly Dominated Strategies; and it is definable inside \mathcal{M} by means of a formula of a modal μ -calculus with inflationary fixed-points. The same for iterated announcement of *SR* and game-theoretic Rationalizability.¹⁶⁶

9.12 Belief, update and revision in extensive games

So far, we have considered players' knowledge in games, which came in various kinds, past observation-based and future-oriented procedural. Next, we consider the concomitant notion of belief, which can also have observational and procedural aspects. To do so, we return to the models in Sections 9.7 through 9.9. Level-One models will add doxastic structure using plausibility relations inside epistemic equivalence classes on the extensive game tree, just as we did in Chapter 6 for belief revision. Intuitively, as we did with pure knowledge, a game tree annotated with players' beliefs in this way may be viewed as a record of stepwise events of knowledge update and belief revision that take place as the game is played, and we will return to this representational issue in Chapter 11.

For the moment, we elaborate another perspective, viz. the Level-Two view of branching trees, a vivid way of modeling beliefs about the future which can be related to our earlier belief update principles in several ways. Recall our earlier epistemic-temporal models for branching time. We already saw how knowledge can be modeled, eliminating future histories by public announcement of the current move, resulting in a recursion axiom

$$F_a \diamond \phi \leftrightarrow (F_a T \ \& \ \diamond F_a \phi)$$

This perspective is found in van Benthem 2004, where it is also extended to belief.

Beliefs over time We now add binary relations $\leq_{i,s}$ of state-dependent *relative plausibility* between histories to branching temporal models. As in Chapter 6, we then get a doxastic modality (existential here, for convenience), with both absolute and conditional versions:

¹⁶⁶ As we note in Chapter 14, if the iterated assertion A has 'existential-positive' syntactic form (for instance, *SR* does), the relevant definition can be formulated in the standard μ -calculus.

Definition Absolute and conditional belief.

We set $\mathbf{M}, h, s \models \langle B, i \rangle \phi$ iff $\mathbf{M}, h', s \models \phi$ for some history h' coinciding with h up to stage s and most plausible for i according to the given relation $\leq_{i,s}$. As an extension, $\mathbf{M}, h, s \models \langle B, i \rangle^\psi \phi$ iff $\mathbf{M}, h', s \models \phi$ for some history h' most plausible for i according to the given $\leq_{i,s}$ among all histories coinciding with h up to stage s and satisfying $\mathbf{M}, h', s \models \psi$. ■

For convenience, we mostly consider an absolute state-independent version of plausibility. Now, belief revision happens as follows. Suppose we are at node s in the game, and move a is played which is publicly observed. At the earlier-mentioned purely epistemic level, this event just eliminates some histories from the current set. But there is now also belief revision, as we move to a new plausibility relation \leq_{i,s^a} describing the updated beliefs.

Hard belief update First, assume that plausibility relations are not node-dependent, making them global. In that case, we have precisely belief revision under *hard information*, as already analyzed in Chapter 6. The new plausibility relation is the old one, restricted to a smaller set of histories. Here is the characteristic recursion law which governs this. To understand it, recall that the temporal operator $F_a\varphi$ says that a is the next event on the current branch, and that φ is true on the current branch immediately after a has taken place:

Fact The following temporal principles are valid for hard revision along a tree:

$$F_a \langle B, i \rangle \phi \leftrightarrow (F_a T \ \& \ \langle B, i \rangle (F_a T, F_a \phi))$$

$$F_a \langle B, i \rangle^\psi \phi \leftrightarrow (F_a T \ \& \ \langle B, i \rangle (F_a \psi, F_a \phi))$$

Again, this is analogous to the dynamic doxastic recursion axioms for hard information in Chapter 6. Behind the Fact lies a general assumption that the plausibility order has only been restricted, not changed as a move is played. More abstractly, in the state-dependent version, this may be viewed as a principle of *Coherence* which says that the most plausible histories for i at h', s' are the *intersection* of B_i with all continuations of s' .¹⁶⁷

Soft update So far, our Level-Two belief dynamics is only driven by hard information, as moves are observed. What would be a game analogue for events of soft information and

¹⁶⁷ Similar principles have been stated in Bonanno 2007 which formalizes *AGM* theory.

matching plausibility changes among worlds, which were the more exciting theme in Chapter 6? This richer setting seems to call for two additional features:

- (a) models with the above state-dependent plausibility relations, so that a move to a next state can change the plausibility pattern on the remaining histories,
- (b) events beyond mere observable moves serving as soft triggers – say, additional signals or ‘hunches’ that players receive in the course of a game.

This time, recursion axioms will not be ‘absolute’ like the preceding ones, but they will reflect policies for plausibility change. We will not pursue this topic here – but in Chapter 11, we do consider the related issue of belief revision under soft information when the plausibility relation runs, not between whole histories, but nodes of the branching tree.

More local dynamics What would be most in line with the spirit of this book, however, are *concrete stepwise procedures* creating beliefs and preferences as we proceed along an extensive game tree. An interesting recent approach of this sort is given in Baltag, Smets & Zvesper 2008, which use assertions of ‘dynamic rationality’ to define those epistemic-doxastic models where plausibility relations for players encode the Backward Induction solution – with ‘intended moves’ indicated by strategies at most plausible worlds. This analysis is not yet quite what we are looking for here, since it works over global Level-Three models encoding all game behaviour beforehand in the worlds, while there are no genuine events of soft update corresponding to real belief revision as the game proceeds.

9.13 Further entanglements: dynamics of rationalization

Finally, as we have noticed several times so far, knowledge and belief are entangled notions, and the entanglement even extends to preference (cf. Chapter 8). Many issues in applying or interpreting games reflect this situation, starting with the solution procedure of Backward Induction itself. We will discuss a few more, just to loosen up the reader’s mind as to what should eventually be described by the logics of this chapter. Here are a few dynamic scenarios from van Benthem 2007 extending well-known facts about games.

Observed actions and rationalization Rationality in the sense of decision theory and game theory often has its greatest successes, not in predicting behaviour in advance, but in

rationalizing it afterwards. We do this all the time. If your preferences between outcomes of some game are not known, I can always *ascribe* a preference which makes your action rational, by assuming you prefer the outcome of the chosen action over the others. This style of rationalization carries over to more complex interactive settings. But now one must also think about *Me*, i.e., the player you are interacting with. The following seems folklore.

Let a finite two-player extensive game G specify my preferences, but not yours. Moreover, let our strategies $\sigma_{me}, \sigma_{you}$ for playing G be fixed in advance, yielding an expanded game model M . All this is supposed to be common knowledge.¹⁶⁸ When can we rationalize your given behaviour σ_{you} to make $\sigma_{me}, \sigma_{you}$ the Backward Induction (*'BI'*) solution of the game? To achieve this, we have total freedom to set your preferences between outcomes of the game, independently from my given ones. Even so, not all models M of this sort support rationalization toward Backward Induction. My given actions encoded in σ_{me} must have a certain quality to begin with, related to my preferences. Note that, at any node of the game tree, playing on according to our given strategies fixes a unique outcome of the game. What is clearly necessary for any successful *BI*-style analysis is the following version of the rationality assertion that we used for analyzing Backward Induction in Section 9.9:

Definition Best-responsiveness.

A game is *best-responsive for Me* given two strategies $\sigma_{me}, \sigma_{you}$ if my σ_{me} always chooses a move leading to an outcome which is at least as good for me as any other outcome that might arise by choosing any other move, and then continuing with the given $\sigma_{me}, \sigma_{you}$. ■

Theorem In any finite game model as described that is best-responsive for *Me*, there exists a preference relation for *You* among outcomes that makes the unique path playing our given strategies against each other the Backward Induction solution.

Proof One starts with final choices for players at the bottom of the game tree, assigning values reflecting preferences for you as described. (My choice was 'rational' automatically, given best-responsiveness.) Now proceed inductively. As *my* turns higher up in the tree are best-responsive for *Me*, I am consistent with *BI*, as the strategies in the sub-games after my

¹⁶⁸ We do just a few dynamic scenarios here: the reader will see many options for variations.

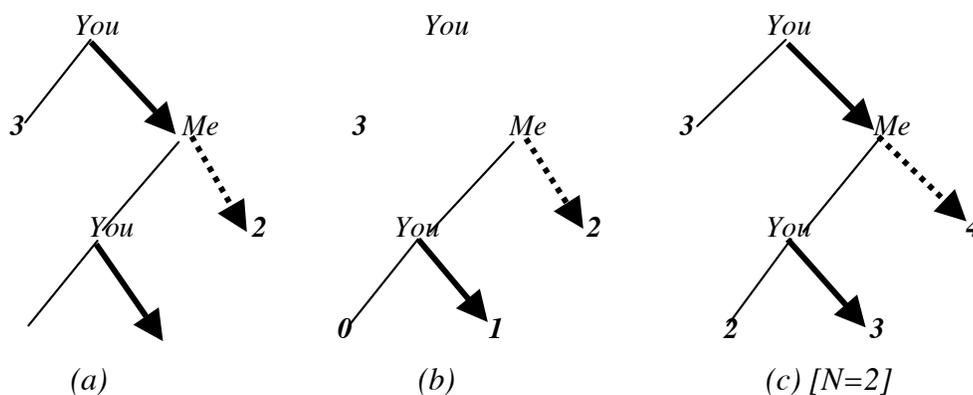
moves are already in accordance with *BI*. Next, let it be *your* turn, with the same inductive assumption. In particular, these sub-games already have *BI*-values for both you and me. Now suppose your prescribed move a in σ_{you} leads to a sub-game G_a with a lower *BI*-value for *You* than some other available move. In that case, a simple trick makes a best for you. Take some fixed number N large enough so that adding N to your utility in all outcomes of G_a makes them better than all outcomes reachable by your other moves than a . Now note:

Raising all your values of outcomes in a game tree by a fixed amount N does not change the *BI*-solution, though it raises your total value by N .

Doing this to a 's subtree, your given move at this turn has become best. ■

Example Rationalizing given strategies.

The procedure runs bottom-up. Bold face arrows mark your given moves, dotted arrows mine. Bold-face numbers at leaves indicate values for outcomes that we postulate for *You*. Think of your value 2, 3 to the left in (b) as *BI*-values assigned by the procedure already – while mine are not indicated (but we can assume my choice are best-responsive):

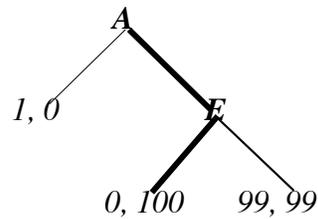


Of course, these numbers can be assigned in many ways to get *BI* right. ■

Given actions and postulated beliefs Here is another dynamic scenario. We just used *BI* to rationalize given strategies in a game G , and implicitly, players' beliefs, postulating preferences for one of the players. Now, we start from given preferences for both players, and *modify the beliefs* of the players to rationalize the given behaviour.

Example Adjusting beliefs with given preferences.

Suppose you choose *Right* in our game with the 'bad outcome' $(1, 0)$:



One can interpret this rationally if we assume you believe I will go *Right* in the next move. This rationalization is not in terms of your preferences, but of your beliefs about me.¹⁶⁹ ■

Again, this readjustment procedure only works under conditions like the rationality used in Section 9.9. Consider a finite game model M with your strategy σ_{you} and your preference relation given (mine does not matter). Suppose *You* can choose move *Left* or *Right*, but all outcomes of *Left* are better for *You* than all those arising after *Right*. Then no beliefs of yours about my subsequent moves can make a choice for *Right* come out ‘best’:

Definition ‘Minimally rational’ game models.

A game model is *minimally rational* for *You* if your strategy σ_{you} never prescribes a move for which each outcome reachable via further play according to σ_{you} and any sequence of moves of *Me* is worse than all outcomes reachable via some other move for *You*.¹⁷⁰ ■

The following result is a counterpart to our *PAL*-style analysis of *BI* in Section 9.9:

Theorem In any game model that is minimally rational for *You*, there is a strategy τ for *Me* against which, if you believe that I will play τ , your given σ_{you} is optimal.

Proof This time, the procedure for finding a rationalizing τ is top-down. Suppose you make a move a right now according to your strategy. Since σ_{you} is minimally rational for you, each alternative move b of yours must have at least one reachable outcome y (via σ_{you} plus some suitable sequence of moves for me) which is majorized by some reachable outcome x via a . In particular, your *maximum value* reachable by playing a always majorizes some value in each sub-game for another move. Now we describe a strategy for *Me* which makes your given move a optimal. Choose later moves for *Me* in the subgame for a which lead to

¹⁶⁹ Note that this belief-based style of rationalizing need not produce the ‘standard’ *BI* solution.

¹⁷⁰ In case you are the last player to move, this coincides with standard rationality again..

9.14 Conclusion

We have shown how games involve a combination of all the logics that we have discussed so far, including knowledge, belief, and preference. We gave pilot studies rather than grand logical theory, though some of the above themes will be taken up again and elaborated in Chapters 11 and 14. But even at this stage, we hope to have shown that logical dynamics and games is an exciting area for research with both formal structure and intuitive appeal.

9.15 Appendix: further topics and open problems

Designing concrete knowledge games The games studied in this chapter are all abstract. But many epistemic and doxastic scenarios in earlier chapters also suggest concrete games, say, of ‘being the first to know’ certain important facts under constraints in observation and communication. It would be of interest to have such ‘gamifications’ of our dynamic logics.

Game equivalence with preference added Taking preference structure seriously would give many topics in our earlier sections a new twist. For instance, van Benthem 2002 considers the issue ‘when two games are the same’ when preference structure is present, looking at intuitions like ‘supporting *the same strategic equilibria*’. These deviate considerably from earlier notions like modal bisimulation, or ‘power equivalence’.

Preferences between actions? Games also suggest a new issue, which would have made sense in Chapter 8 already, about combining *propositional dynamic logic* and *preference logic*. Since *PDL* has formulas as properties of states and programs as inter-state relations, we can put preference structure at *two levels*. One kind runs between states, interpreting standard modal operators as in Chapter 8, which focused on comparison between worlds. The other locus is preferences between state transitions (‘moves’, ‘events’) as in van der Meijden 1996, van Benthem 1999. Compare the distinction in ethics between ‘deontology’ and ‘consequentialism’. Do we compare the worlds resulting from actions when judging our duties, or do we qualify those actions themselves as ‘better’ or ‘obligatory’? Games suggest a preference logic directly on moves, and our guess is that this can be developed in a manner analogous to our dynamic-epistemic logic *DEL* in Chapter 5 with product update for events, perhaps using rules more like the priority update of Chapter 6.

Knowledge that and know-how Taking strategies seriously also raises new issues, such as the famous distinction between ‘knowing that’ and ‘knowing how’. Strategies are ways of achieving goals, and hence they represent *procedural know-how*. What does it mean to ‘know a strategy’ in epistemic dynamic logics, and can we develop significant logics of strategies in games with imperfect information? Chapter 13 studies some of these issues, including *epistemically transparent* strategies σ with the property that, when part of σ has been played, players still know that playing the rest will achieve the intended result.

Local dynamics inside given games We already mentioned the lack of a local marriage of *DEL* style events and moves in extensive games, tighter than what we have provided here.

Soft update in games Our dynamic scenarios involved mostly public observation of moves. What natural extension of games would support the soft information events of Chapter 6?

Agent diversity and memory We characterized the *DEL*-like games of imperfect information in terms of players with Perfect Recall. But there are many others, such as the scenario of the ‘Drunken Driver’. Can the extended dynamic logics that allow for agents with *bounded memory* (Chapters 3, 4) analyze such further game-theoretic scenarios?

Parallel actions Many games involve simultaneous moves by players, but so did epistemic scenarios of speaking simultaneously. How can we introduce parallelism into both?

From finite to infinite games All examples in this chapter are finite games. Can we extend our dynamic logic analysis to deal with infinite games, such as repeated Prisoner’s Dilemma, and the games normally dealt with in evolutionary game theory (Hofbauer & Sigmund 1998)? This transition to the infinite was implicit in the iterated announcement scenarios of Chapter 3, and it is also natural in modal process theories of computation (Bradfield & Stirling 2006). Taking this view leads naturally to the epistemic and doxastic temporal logics of Chapter 11, where we will show more precisely how *DEL* fits into ‘Grand Stage’ views of temporal evolution of action, knowledge, and belief.

